

**AIM WORKSHOP ON
COMPUTABLE STABILITY THEORY
PROBLEM SESSION**

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1. QUESTIONS

1.1. Computable approximability (Calvert).

Definition. \mathcal{M} is **computably approximable** if for every computable $\mathcal{L}_{\omega_1, \omega}$ sentence φ true of \mathcal{M} there is a computable $\mathcal{N} \models \varphi$ such that $\text{SR}(\mathcal{N}) < \omega_1^{\text{CK}}$, where SR denotes Scott rank.

Question 1. *Is it the case that every computable structure is computably approximable?*

Date: August 12, 2013.

Question 2. Let $\varphi \in \mathcal{L}_{\omega_1, \omega}$ be a satisfiable sentence of quantifier rank α , and suppose that either $\omega_1 > \beta > \alpha$ or $\beta > \alpha + \omega$. Is there an $\mathcal{M} \models \varphi$ such that $\text{SR}(\mathcal{M}) < \beta$?

1.2. Strongly minimal nontrivial locally modular nonorthogonal groups (Medvedev).

Let \mathcal{M} be a strongly minimal nontrivial locally modular structure. There is a nonorthogonal interpretable strongly minimal group G .

Question 3. How difficult is it to find a presentation of G in terms of \mathcal{M} ?

Question 4. How difficult is it to find a presentation of \mathcal{M} in terms of G ?

Consider a three-to-one map $(\mathbb{Q}, +) \leftarrow (\mathcal{M}, \oplus)$ where (\mathcal{M}, \oplus) is strongly minimal.

Question 5. Must \oplus be definable in some $\mathbb{Q} \times F$, where F is a finite set?

1.3. Continuous sections (Miller).

Consider T stable. Then for all $\mathcal{M} \prec \mathcal{N} \models T$, the map $S_1(\mathcal{N}) \rightarrow S_1(\mathcal{M})$ has a continuous section, which sends p to the unique non-forking extension.

Question 6. Is there a computable section?

Question 7. How complicated is the map $\varphi \mapsto d_p\varphi$ (possibly with uniformity in p)?

Question 8. Does the existence of a computable section give other computable information for other characterizations of stability?

1.4. κ^+ -computable categoricity (Knight).

Definition. Let κ be a cardinal. Then a set is κ^+ -recursively enumerable when it is Σ_1 on L_{κ^+} .

Definition. \mathcal{K} is relatively κ^+ -categorical when for any two κ^+ -computable \mathcal{N}, \mathcal{M} of cardinality κ^+ , they are isomorphic in $L_{\kappa^+}(\mathcal{M}, \mathcal{N})$.

Recall that when an AEC is quasiminimal excellent, it is κ -categorical for all uncountable κ .

Question 9. *Let \mathcal{K} be a quasiminimal excellent class, and $\lambda < \kappa$. Suppose that \mathcal{K} is κ^+ -computably categorical. Must \mathcal{K} be λ^+ -computably categorical?*

Question 10. *Suppose that \mathcal{K} is relatively κ^+ -computably categorical. Must \mathcal{K} be relatively λ^+ -computably categorical?*

1.5. Σ -definable isomorphisms for copies of \mathbb{C} (Goncharov).

Question 11. *Let $\mathbb{A} = \text{HF}(\mathbb{C})$, and suppose $K \cong \mathbb{C}$ is Σ -definable in \mathbb{A} . Is there a Σ -definable isomorphism?*

The answer is yes if we replace \mathbb{C} by \mathbb{R} and both $(K, \oplus, \odot) \cong \mathbb{R}$ and $K \subseteq \mathbb{R}$ hold. It is open if merely $K \subseteq \mathbb{R}^2$.

1.6. λ -many models of each cardinality $\lambda \geq \aleph_1$ (Greenberg).

Question 12. *(Assume $V = L$ if it makes things easier.) Suppose that for all $\lambda \geq \aleph_1$ a theory T has at most λ -many models of cardinality λ . Must each such model have a λ -computable presentation?*

Such a T is necessarily ω -stable and non-multidimensional.

The question is true if T is \aleph_1 -categorical.

Question 13. *What if T has finitely many models in \aleph_1 ? (Maybe look at models in \aleph_n .)*

1.7. **Non-abelian free groups (Knight).** Consider n -generated groups with a single relator of length at most t . (For each t, n there are finitely many such groups.) For every sentence φ define

$$h_{n,t,\varphi} = \frac{|\{G \in H_{n,t} : G \models \varphi\}|}{|H_{n,t}|}.$$

Conjecture 14. *The limit $\lim_{t \rightarrow \infty} h_{n,t,\varphi}$ exists, and always takes the value 0 or 1, moreover in a way that may depend on φ but not on n .*

Conjecture 15. *Furthermore, the asymptotically almost sure (a.a.s.) theory determined by this 0 – 1 law is that of \mathbb{F}_2 .*

1.8. Standard systems of RCF (Marker).

Question 16. *What are the possible standard systems of recursively saturated real closed fields? (This is essentially asking: What are the possible sets of cuts of \mathbb{Q} that are realized in some models?)*

The ideal answer might be “all Scott sets”.

1.9. Models of \aleph_1 -categorical theories. (Andrews).

Conjecture 17. *For any \aleph_1 -categorical T there is an n such that if T has a computable model then every countable model has a presentation computable in $0^{(n)}$.*

Note that if T is strongly minimal then $n = 4$ works.

1.10. Computable prime models (Andrews).

Question 18 (Millar). *Let T be a decidable theory having countably many countable models. When must the prime model have a decidable presentation? (Note that ω -stability suffices.)*

Question 19. *Let T be a decidable theory having countably many countable models. What do we need to know to build a computable prime model of T ? (Of course ω -stability again suffices.)*

1.11. Turing degrees of DCFs (Calvert).

Question 20 (Harizanov). *Let \mathbf{d} be a Turing degree. Is there a differentially closed field with a copy that is computable in \mathbf{d} and such that every copy computes \mathbf{d} ?*

1.12. Spectrum of totally categorical theories (Andrews).

Definition. $\text{Spec}(T) = \{d : \text{there is a model of } T \text{ computable in } d\}$.

Conjecture 21. *If T is totally categorical then $\text{Spec}(T)$ is a cone.*

This is true in a finite language.

1.13. Model theoretic consequences of Erdős-Rado (Greenberg).

Task 22 (Hirschfeldt). *Find proofs in second-order arithmetic of model-theoretic consequences of Erdős-Rado (e.g., forking = dividing in simple theories).*

1.14. Borel complexity of isomorphism (Marker).

Question 23. *Let T' be an expansion of T by one constant. Suppose $\cong_{T'}$ is Borel complete. Is \cong_T Borel complete?*

Question 24. *Suppose \cong_T is Borel complete. Is $\cong_{T'}$ Borel complete?*